

# Counterfactuals Modulo Temporal Logics

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# Counterfactuals

*“If  $\varphi$  had been true, then  $\psi$  would have been true, too.”*

$$\varphi \square \rightarrow \psi$$



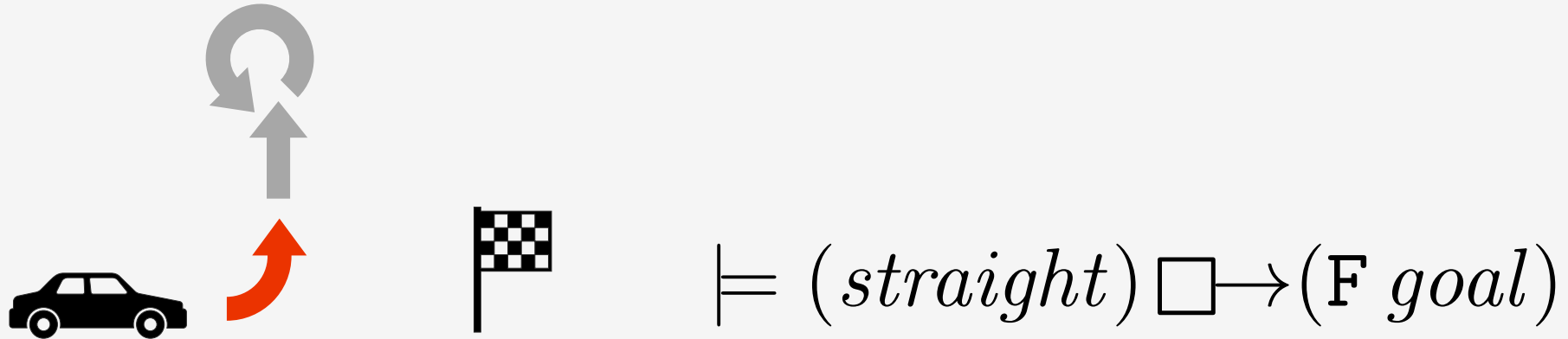
# Counterfactuals Modulo Temporal Logics

“If the car had **moved straight** at the first time point, then it would have **reached its goal eventually**.”

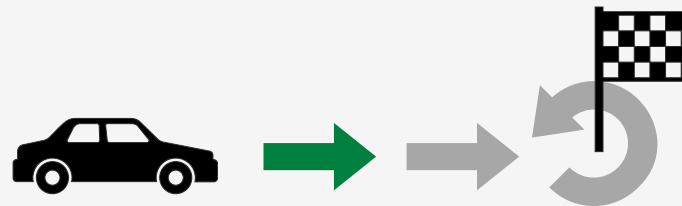
$$\left( \textit{straight} \right) \square \rightarrow \left( \mathbf{F} \textit{goal} \right)$$



# Counterfactuals = Variably Strict Conditionals

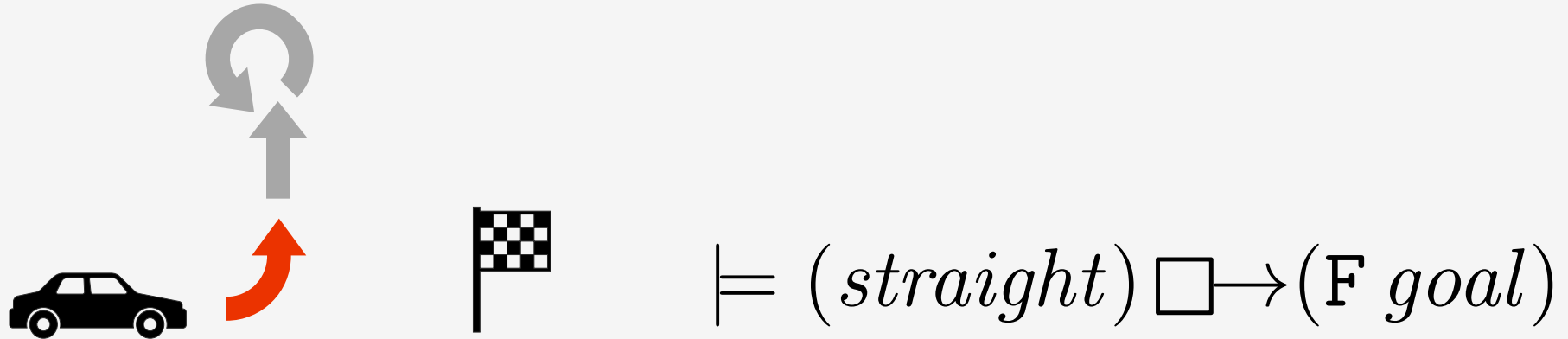


because the closest counterfactual world is:

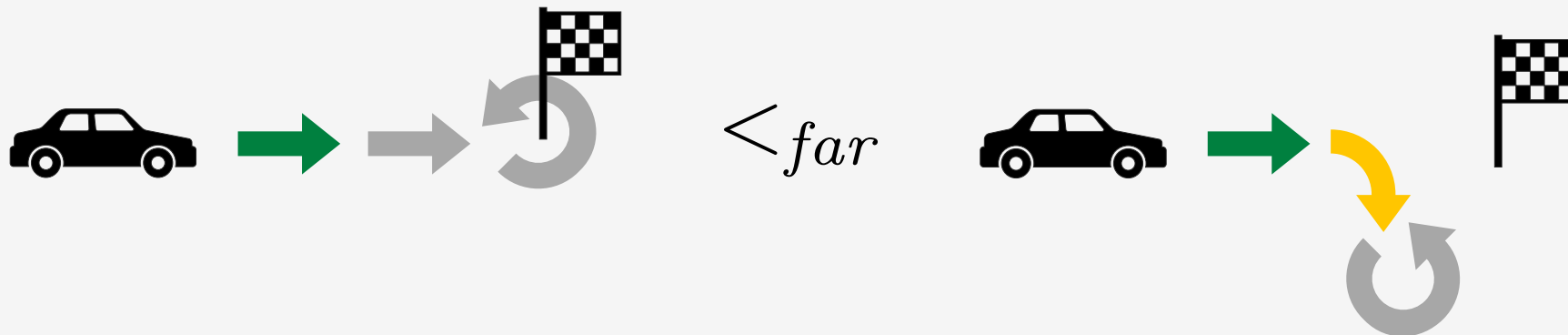




# Counterfactuals = Variably Strict Conditionals



counterfactual worlds further away do not matter:





# Applications of Counterfactual Reasoning



Analyzing **causality** [Halpern '15], [Leitner-Fischer '15], [Coenen et al. '22]



Generating **explanations** for, e.g., model checkers [Beer et al. '09], [Wachter et al. '18]



Counterfactual **fairness** [Kusner et al. '17]



# Applications of Counterfactual Reasoning



Analyzing **causality** [Halpern '15], [Leitner-Fischer '15], [Coenen et al. '22]

**Definition 5 (Property Causality).** Let  $\mathcal{T}$  be a system,  $\pi \in \text{traces}(\mathcal{T})$  a trace,  $C \subseteq (2^I)^\omega$  a cause property, and  $E \subseteq (2^O)^\omega$  an effect property. We say that  $C$  is a cause of  $E$  on  $\pi$  in  $\mathcal{T}$  if the following three conditions hold:

**PC1:**  $\pi \models C$  and  $\pi \models E$ , i.e., cause property and effect property are satisfied by the actual trace.

**PC2:** For every counterfactual input sequence  $\sigma \in V_\pi^C$ , there is some contingency  $\pi' \in C_\pi^\sigma$  s.t.  $\pi' \not\models E$ , i.e., the counterfactual trace under contingency does not satisfy the effect property.

**PC3:** There is no  $C'$  s.t.  $C' \subset C$  and  $C'$  satisfies PC1 and PC2.

[Coenen et al. '22]

**Definition 3.1:** (Actual cause)  $\vec{X} = \vec{x}$  is an *actual cause* of  $\varphi$  in  $(M, \vec{u})$  if the following three conditions hold:

AC1.  $(M, \vec{u}) \models (\vec{X} = \vec{x}) \wedge \varphi$ . (That is, both  $\vec{X} = \vec{x}$  and  $\varphi$  are true in the actual world.)

AC2. There exists a partition  $(\vec{Z}, \vec{W})$  of  $\mathcal{V}$  with  $\vec{X} \subseteq \vec{Z}$  and some setting  $(\vec{x}', \vec{w}')$  of the variables in  $(\vec{X}, \vec{W})$  such that if  $(M, \vec{u}) \models Z = z^*$  for  $Z \in \vec{Z}$ , then

(a)  $(M, \vec{u}) \models [\vec{X} \leftarrow \vec{x}', \vec{W} \leftarrow \vec{w}'] \neg \varphi$ . In words, changing  $(\vec{X}, \vec{W})$  from  $(\vec{x}, \vec{w})$  to  $(\vec{x}', \vec{w}')$  changes  $\varphi$  from true to false,

(b)  $(M, \vec{u}) \models [\vec{X} \leftarrow \vec{x}, \vec{W} \leftarrow \vec{w}', \vec{Z}' \leftarrow \vec{z}^*] \varphi$  for all subsets  $\vec{Z}'$  of  $\vec{Z}$ . In words, setting  $\vec{W}$  to  $\vec{w}'$  should have no effect on  $\varphi$  as long as  $\vec{X}$  is kept at its current value  $\vec{x}$ , even if all the variables in an arbitrary subset of  $\vec{Z}$  are set to their original values in the context  $\vec{u}$ .

AC3.  $\vec{X}$  is minimal; no subset of  $\vec{X}$  satisfies conditions AC1 and AC2. Minimality ensures that only those elements of the conjunction  $\vec{X} = \vec{x}$  that are essential for changing  $\varphi$  in AC2(a) are considered part of a cause; inessential elements are pruned. ■

[Halpern&Pearl '05]



# Applications of Counterfactual Reasoning



Analyzing **causality** [Halpern '15], [Leitner-Fischer '15], [Coenen et al. '22]

$$\varphi \wedge \psi \wedge \left( (\neg\varphi \square \rightarrow_{min} \neg\psi) \vee (\neg\varphi \diamond \rightarrow_{min} \neg\psi) \right)$$

[Coenen et al. '22]

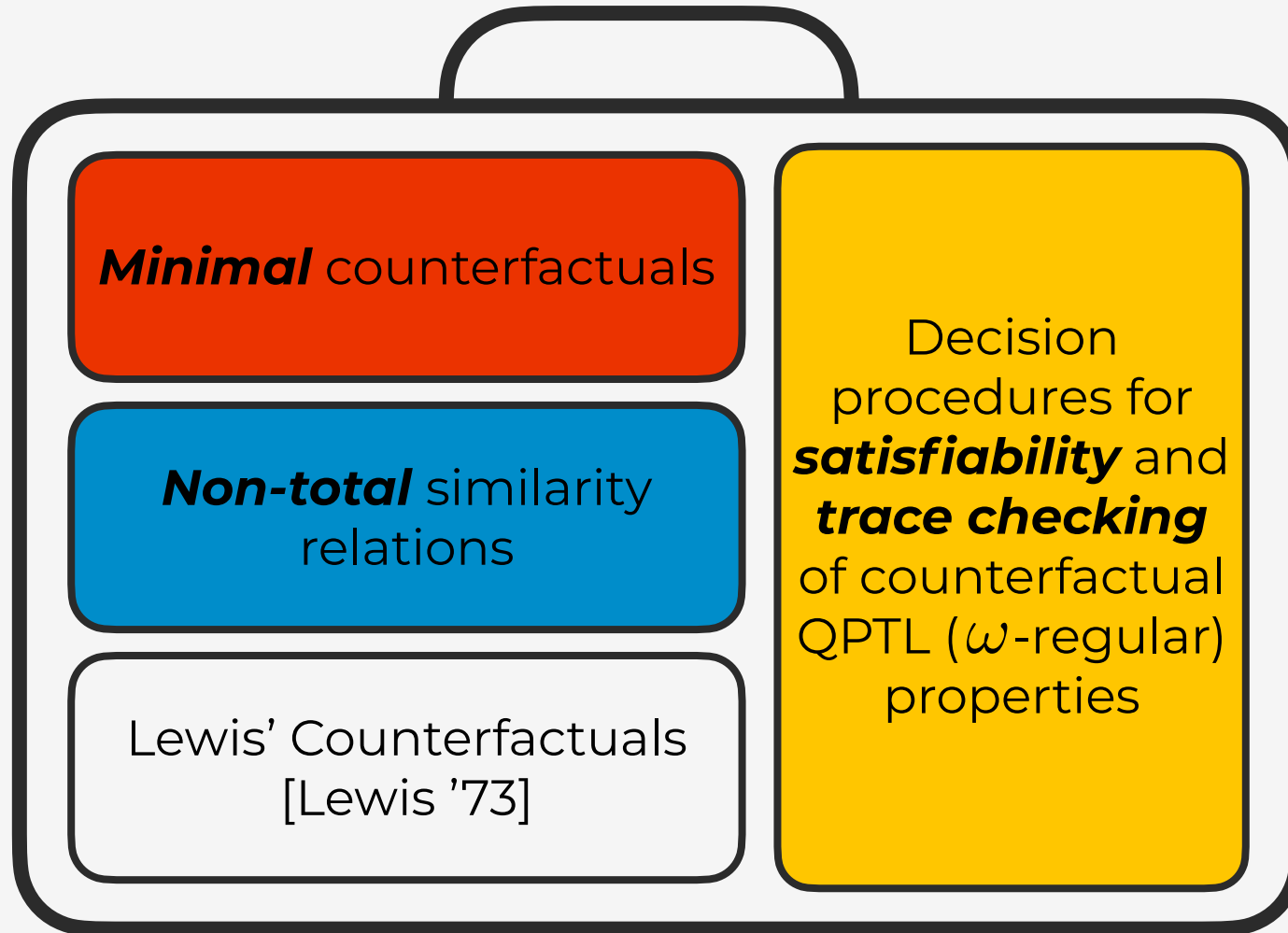
$$\varphi \wedge \psi \wedge (\neg\varphi \diamond \rightarrow_{min} \neg\psi)$$

[Halpern '15]





# Overview



Specification toolbox for, e.g., temporal causality

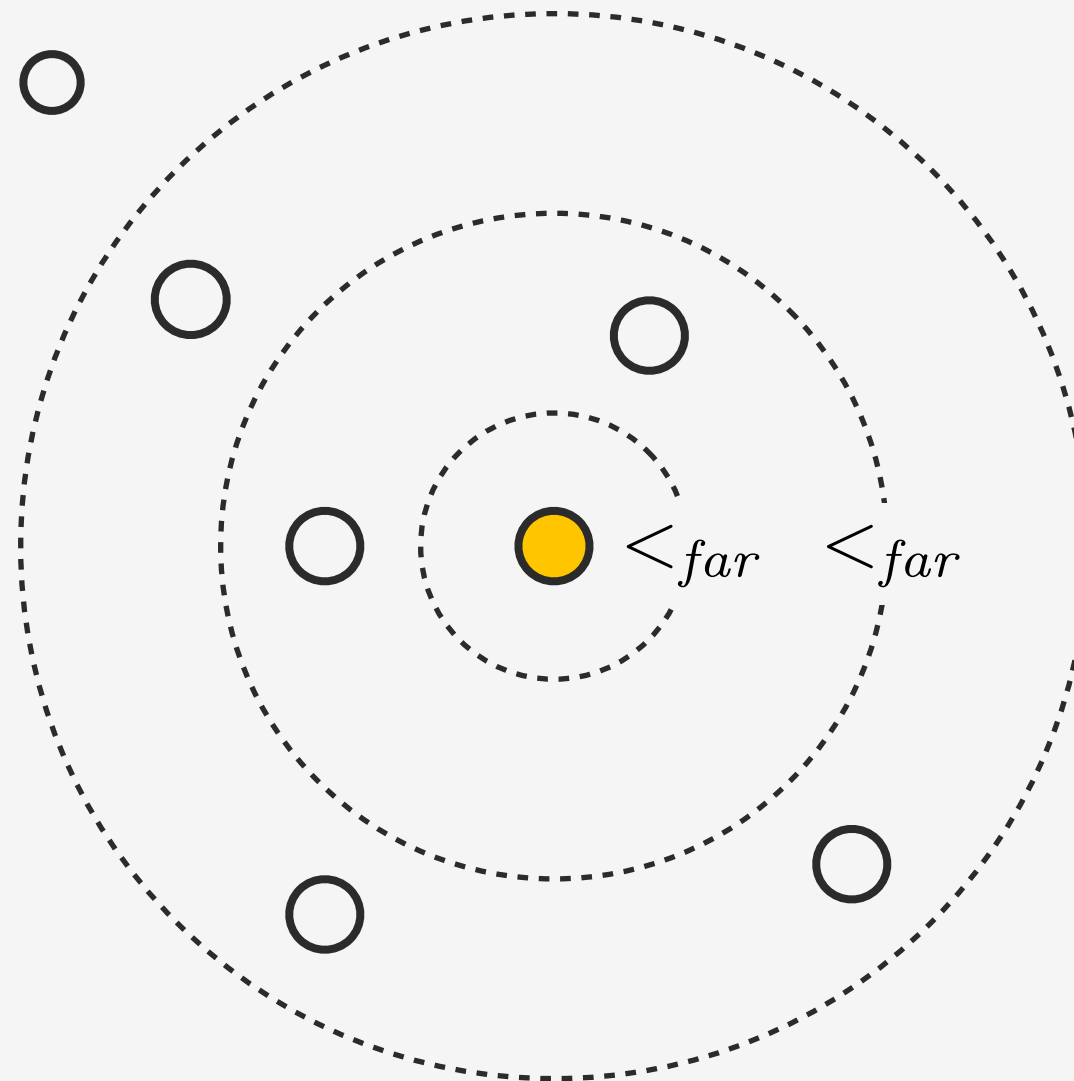


# Similarity Relation $\leq_{far}(\bullet)$

Encodes the distance of a world  $\circ$  from the reference world  $\bullet$ .

Needs to be a total preorder.

$\bullet$  is the unique minimum.





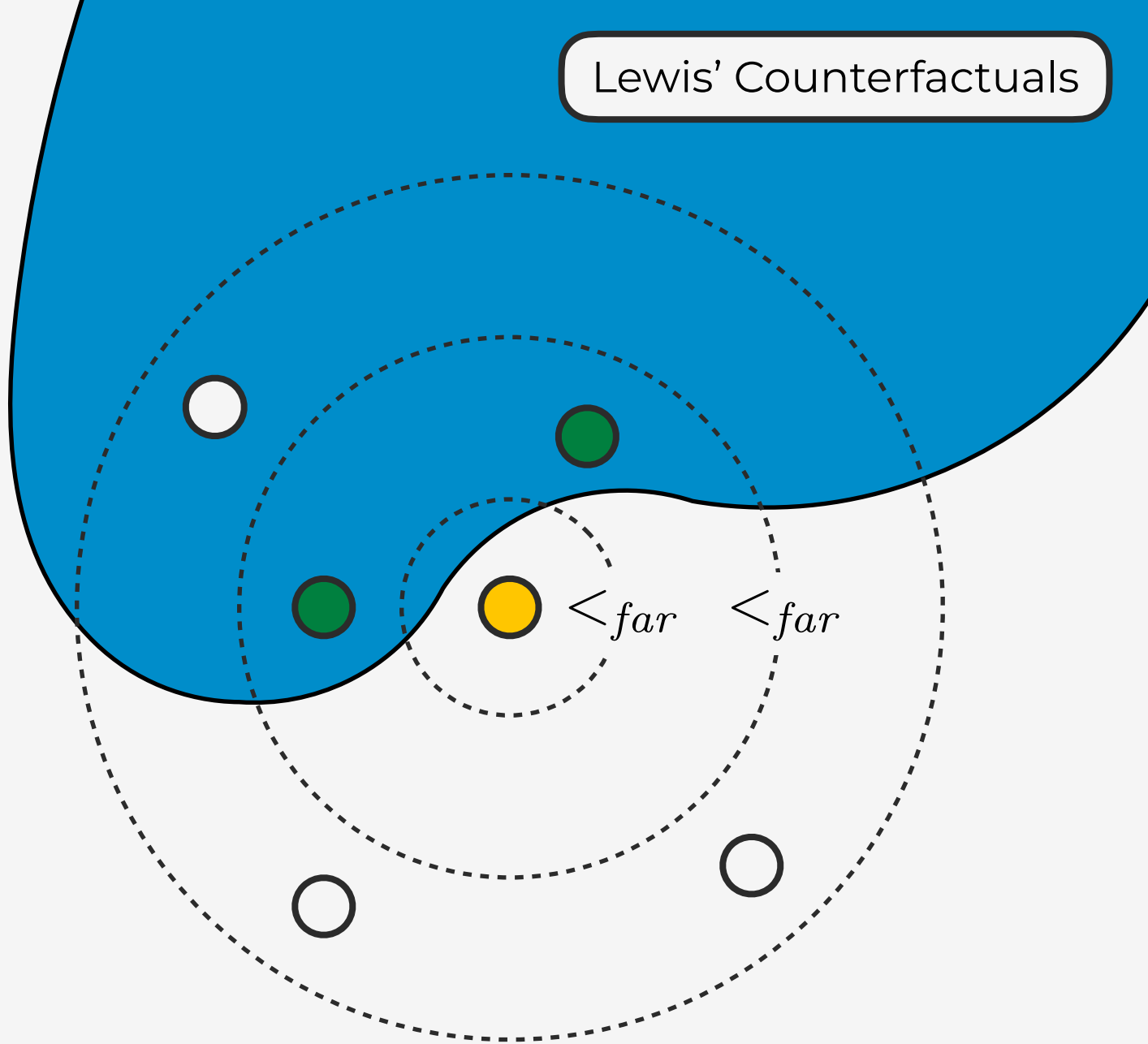
# 'Would' Operator

Lewis' Counterfactuals

$$\varphi \square \rightarrow \psi$$

All closest worlds satisfying  $\varphi$   
have to satisfy  $\psi$ .

Worlds in spheres further  
away do not matter.





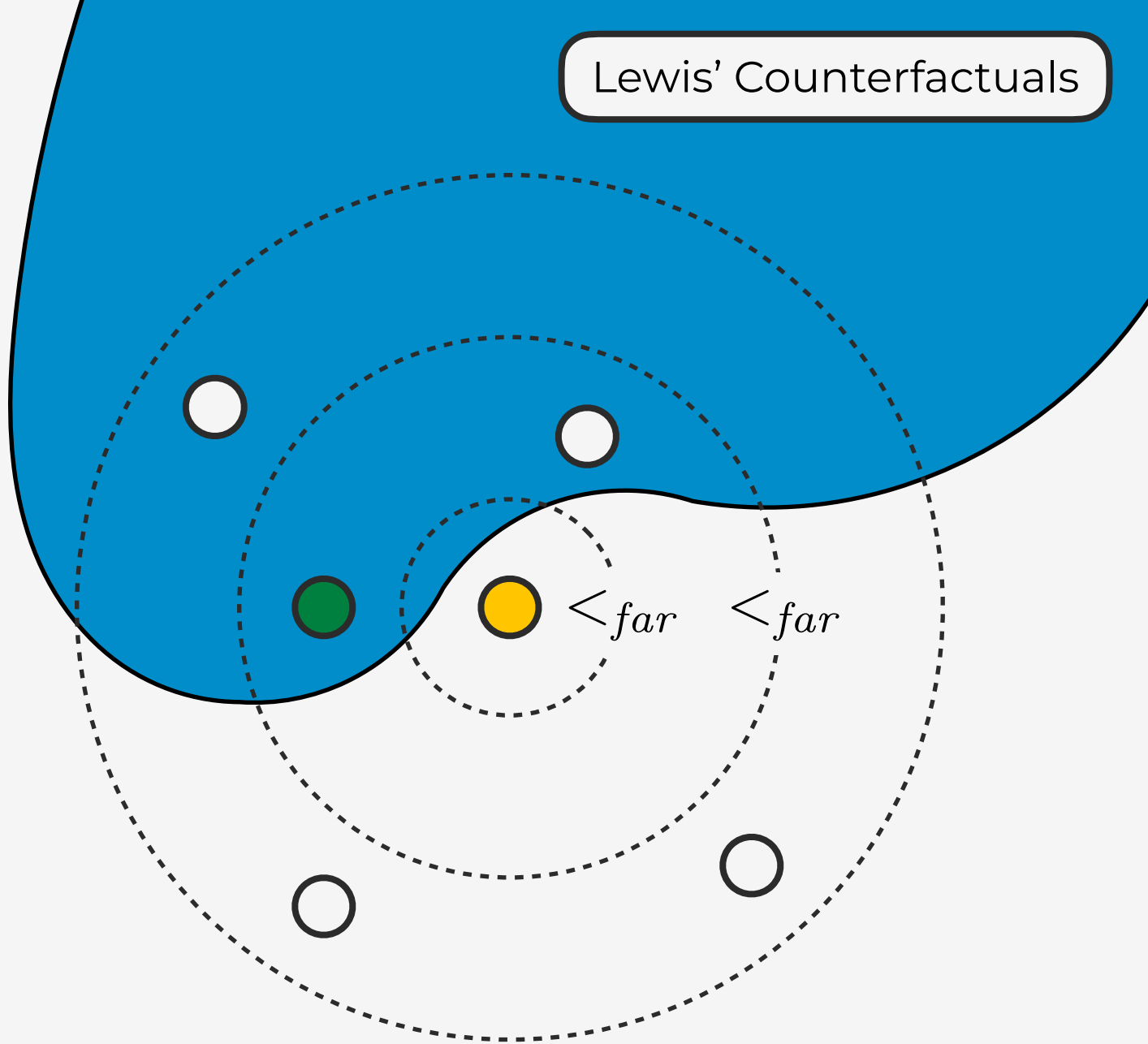
# 'Might' Operator

Lewis' Counterfactuals

$$\varphi \diamond \rightarrow \psi$$

Some closest world satisfying  $\varphi$   
has to satisfy  $\psi$ .

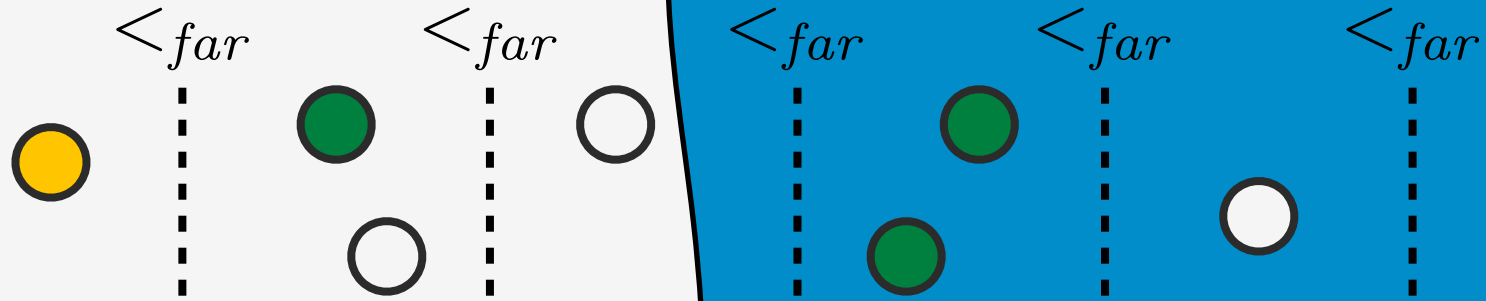
Again, worlds in spheres  
further away do not  
matter.





# The Limit Assumption

*“There always exist well-defined closest worlds satisfying  $\varphi$  .”*

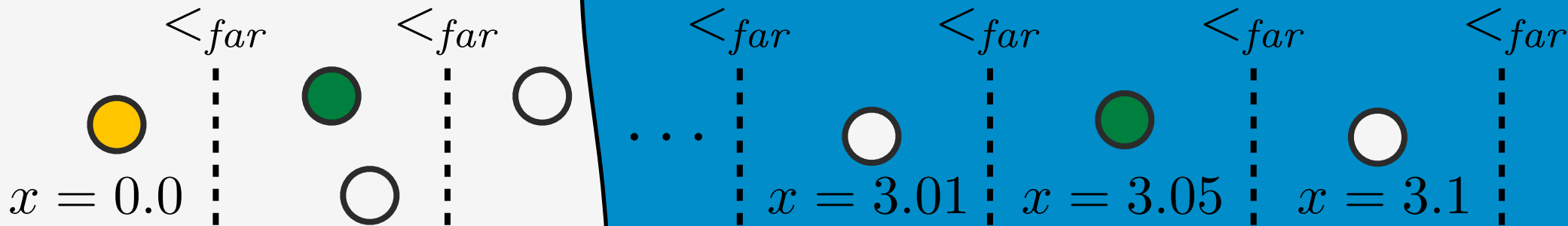




# The Limit Assumption

Lewis' Counterfactuals

*“There always exist well-defined closest worlds satisfying  $\varphi$  .”*



Generally not satisfied and already **rejected by Lewis.**

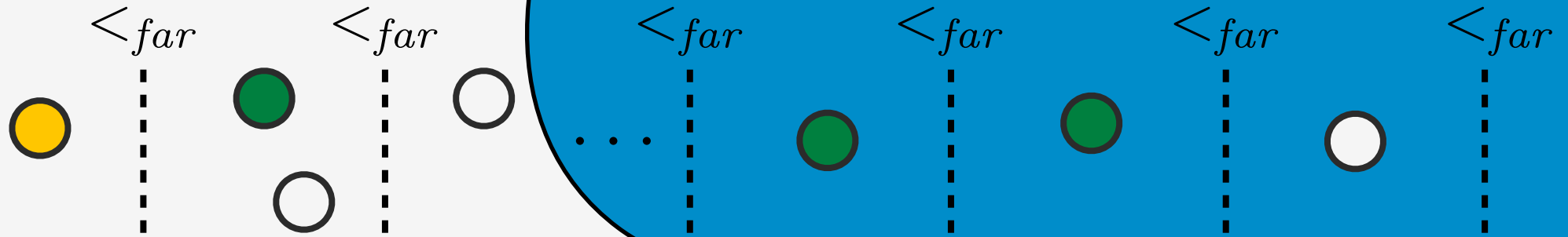
$x \leq 3.0$      $x > 3.0$



# Semantics of 'Would'

Lewis' Counterfactuals

$$\text{Yellow Circle} \models \varphi \square \rightarrow \psi \quad \text{iff:}$$



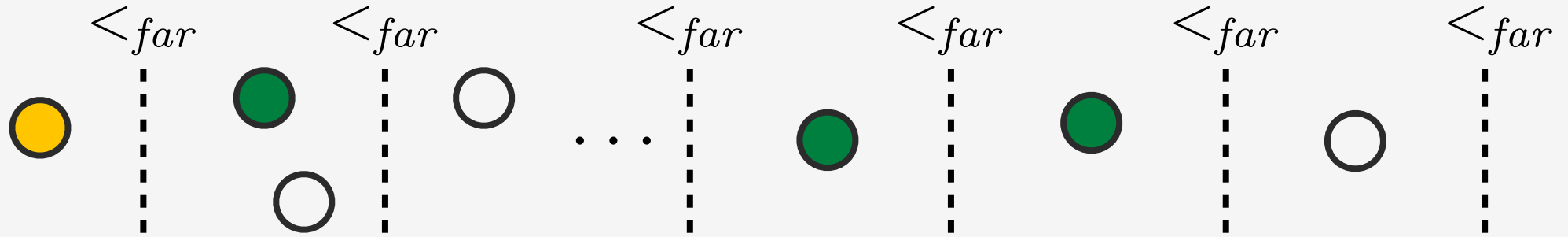
There is a threshold world after which all closer  $\varphi$ -worlds satisfy  $\psi$ .

$$(1) \exists \text{Green Circle} : \text{Green Circle} \models \varphi \wedge \forall \text{White Circle} : \text{White Circle} \leq_{far} \text{Green Circle} \Rightarrow (\text{White Circle} \models \varphi \rightarrow \psi)$$



# Semantics of 'Would'

$$\bigcirc \models \varphi \square \rightarrow \psi \quad \text{iff:}$$



Or: There are no  $\varphi$ -worlds (vacuous case).

$$(1) \text{ or } (2): \forall \bigcirc : \bigcirc \not\models \varphi$$

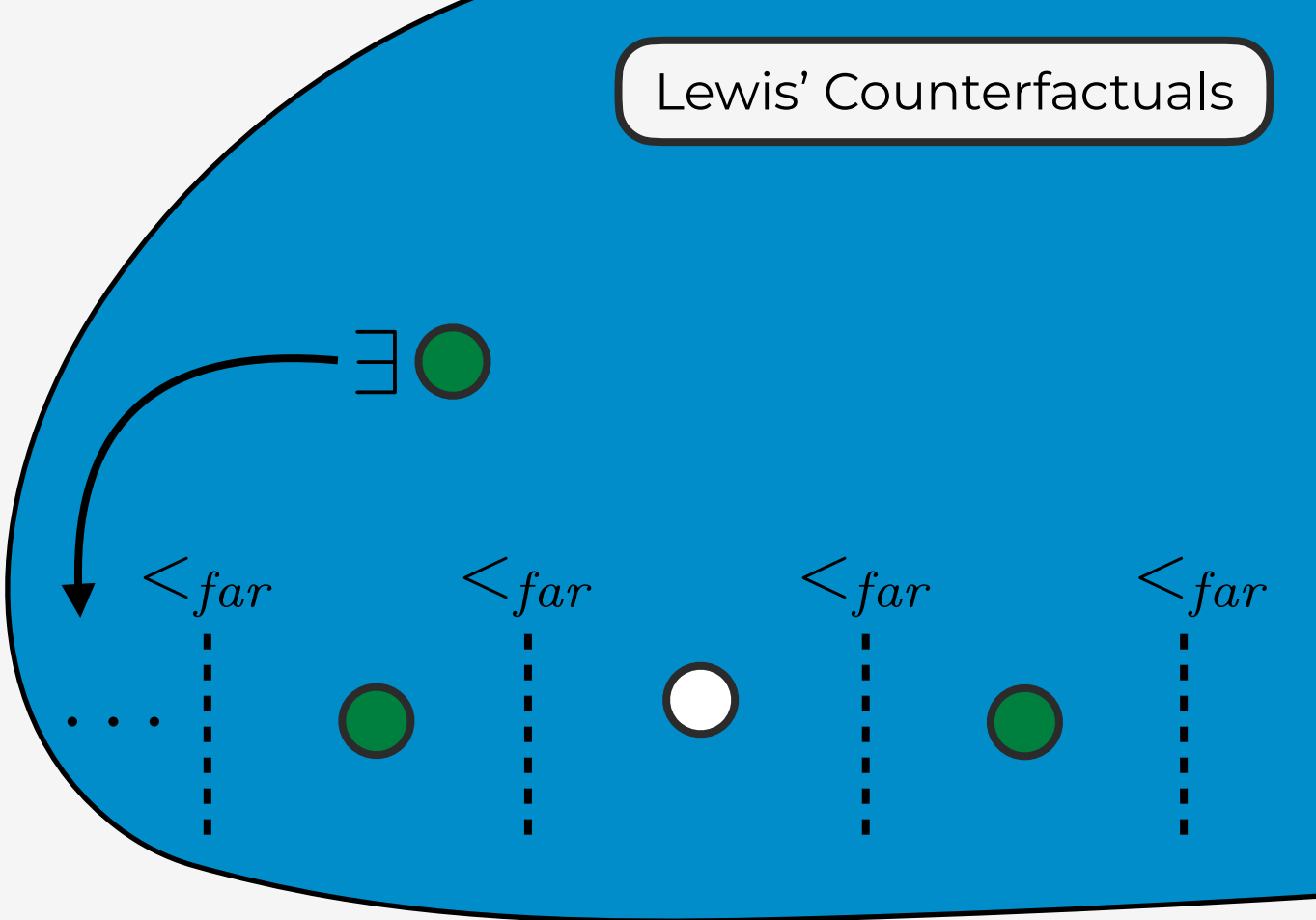
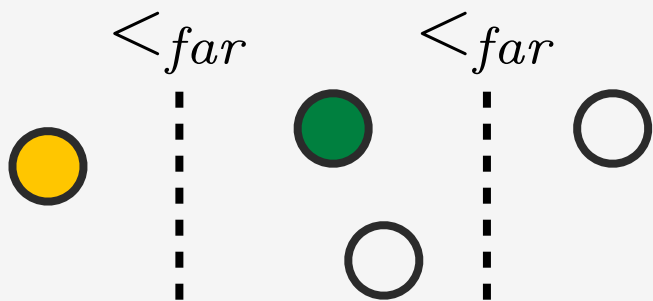




# Semantics of 'Might'

Lewis' Counterfactuals

$$\text{Yellow Circle} \models \varphi \diamond \rightarrow \psi \quad \text{iff:}$$



For any  $\varphi$ -world (and there is at least one) there exists a closer world satisfying  $\psi$ .

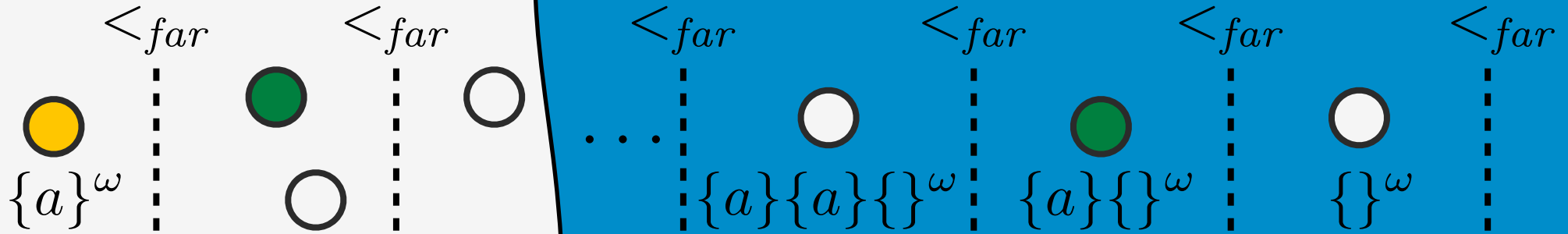
$$\exists \text{White Circle} : \text{White Circle} \models \varphi \wedge \forall \text{Grey Circle} : \text{Grey Circle} \models \varphi \Rightarrow \exists \text{Green Circle} : \text{Green Circle} \leq_{far} \text{Grey Circle} \wedge \text{Green Circle} \models \psi$$



# The Limit Assumption

Lewis' Counterfactuals

“There always exist well-defined closest worlds satisfying  $\varphi$  .”



Also does not hold for temporal properties.

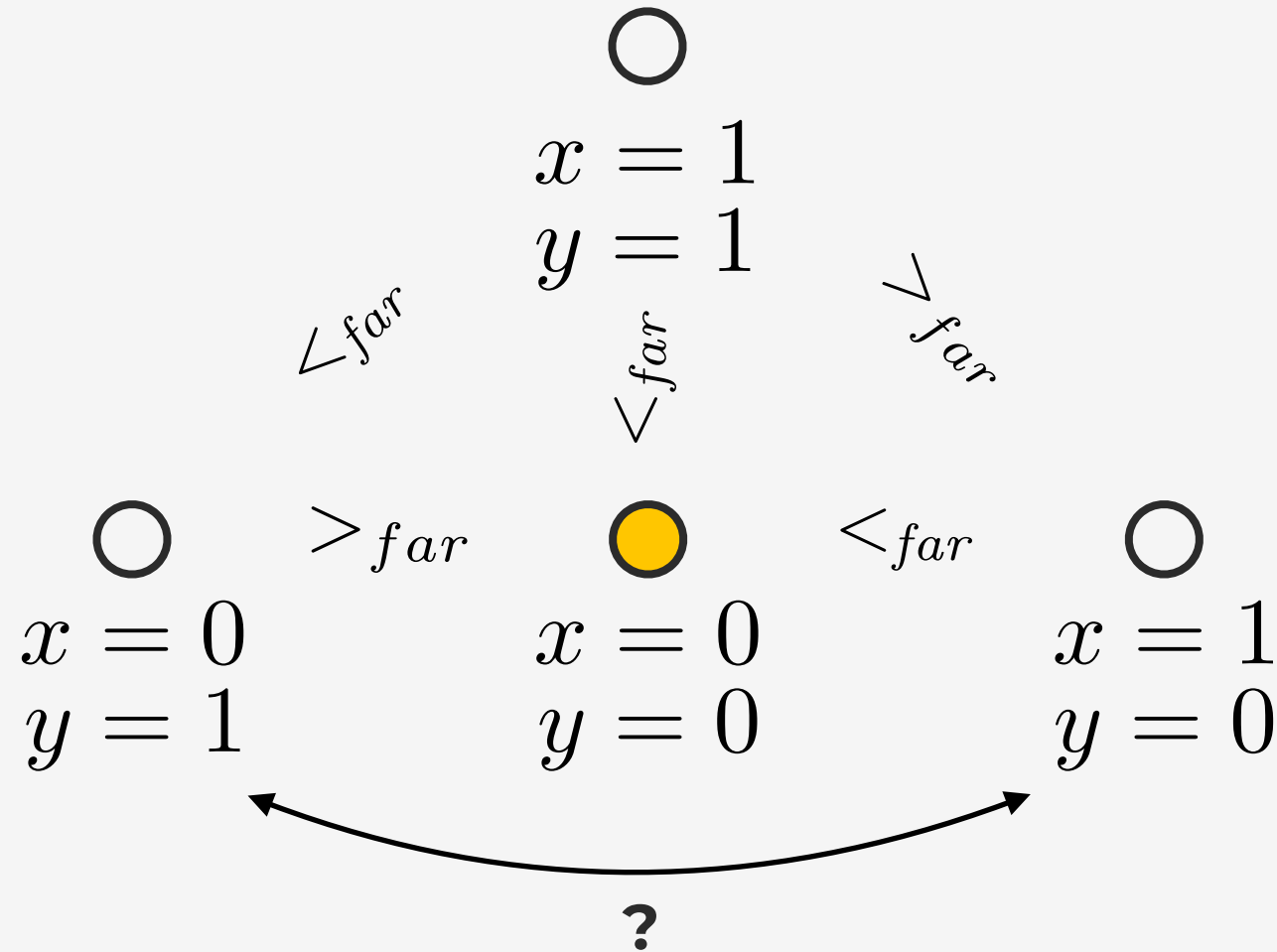
$G F a$

$F G \neg a$



# The Problem with Linearity

Lewis' Counterfactuals

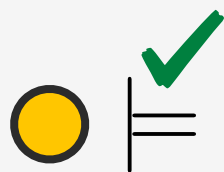


In practice, counterfactual worlds are often incomparable.

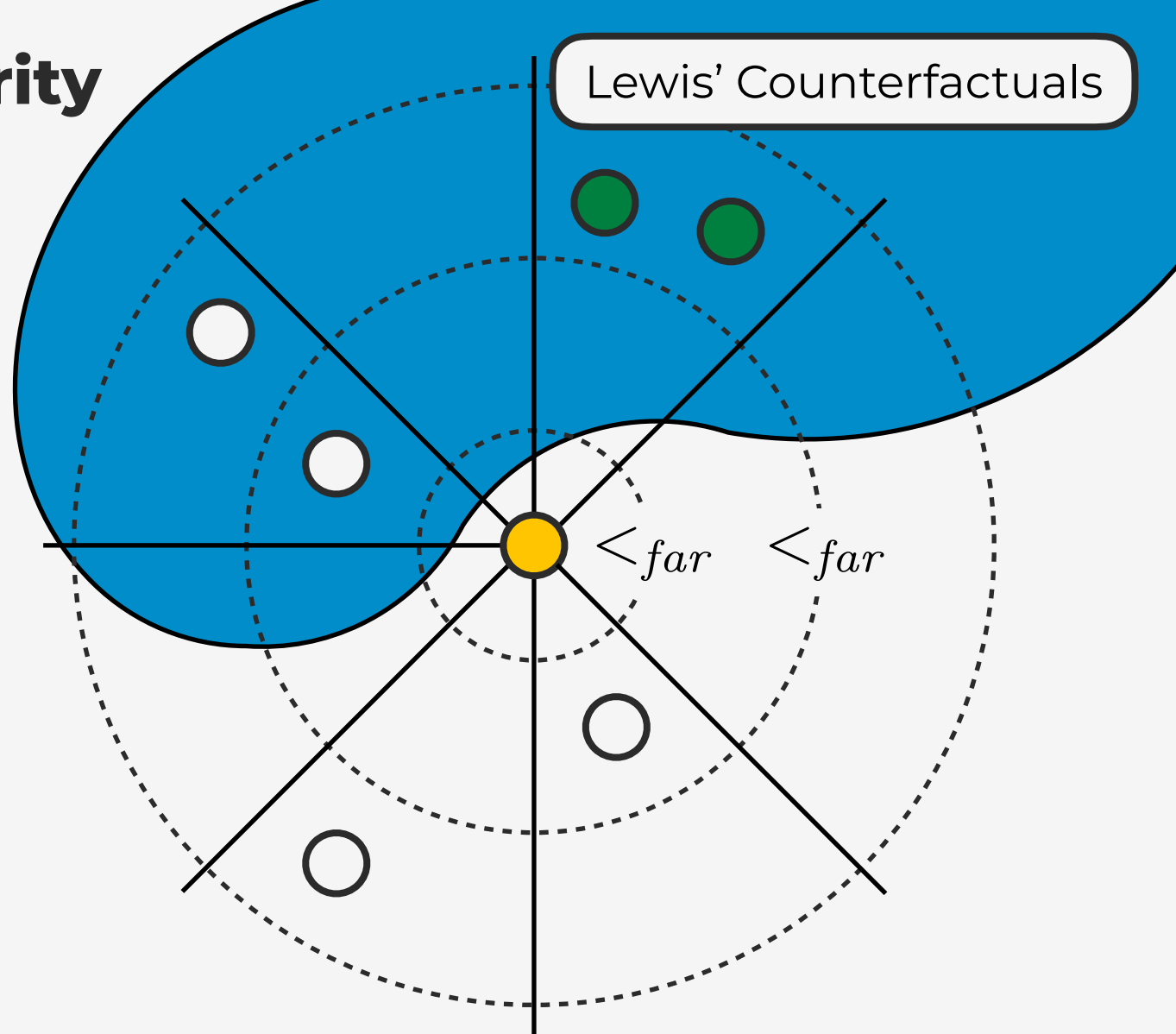


# The Problem with Linearity

Applying Lewis' semantics to non-total orders leads to unintuitive judgements.

  $\models \varphi \square \rightarrow \psi$  iff:

$$(1) \exists \bullet : \bullet \models \varphi \wedge \forall \circ : \circ \leq_{far} \bullet \Rightarrow (\circ \models \varphi \rightarrow \psi)$$

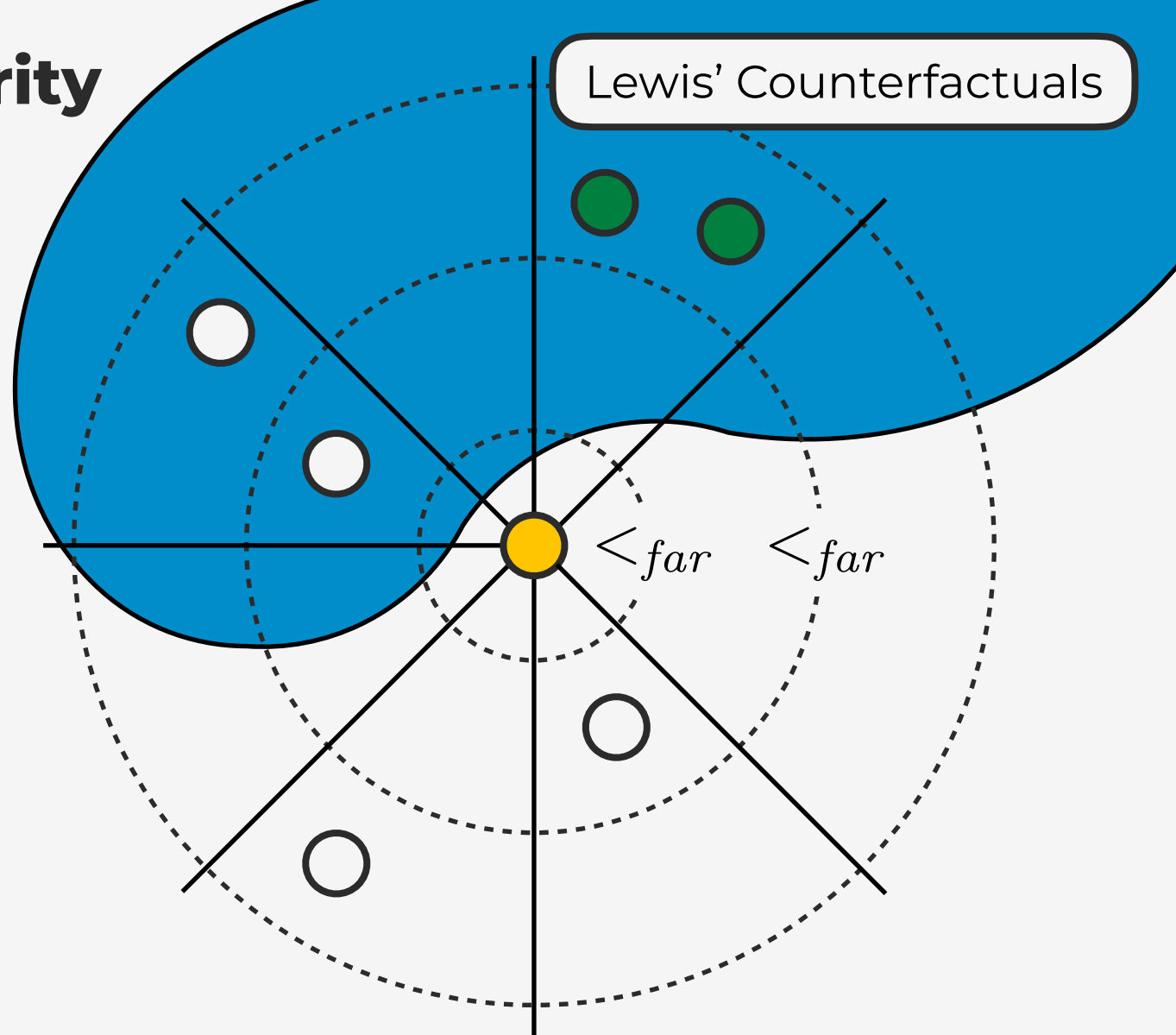




# The Problem with Linearity

Applying Lewis' semantics to non-total orders leads to unintuitive judgements.

$$\text{Yellow Circle} \models \varphi \diamond \rightarrow \psi \text{ iff:}$$




$$\exists \text{ White Circle} : \text{White Circle} \models \varphi \wedge \forall \text{ Grey Circle} : \text{Grey Circle} \models \varphi \Rightarrow \exists \text{ Green Circle} : \text{Green Circle} \leq_{far} \text{Grey Circle} \wedge \text{Green Circle} \models \varphi$$

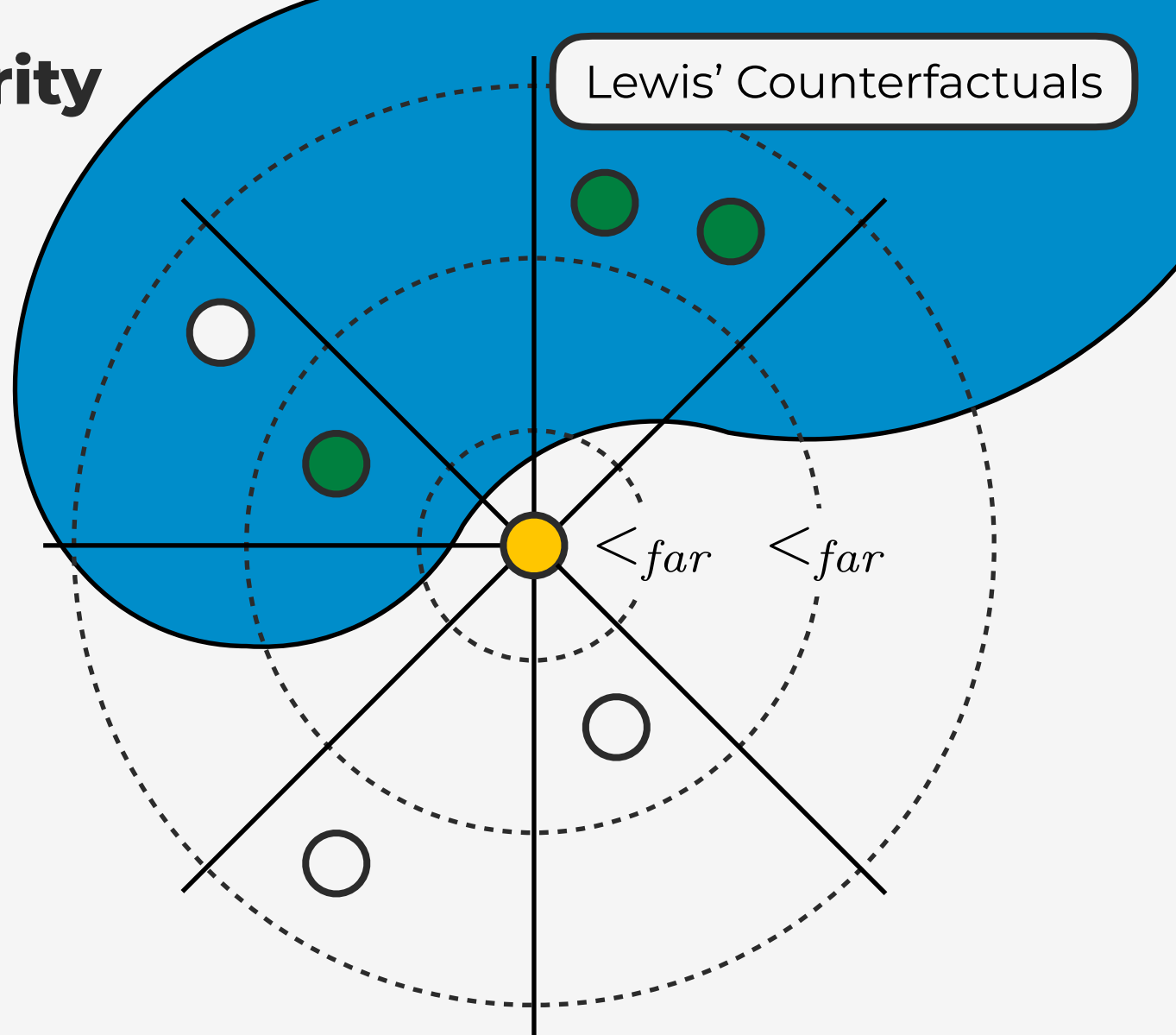


# The Problem with Linearity

Applying Lewis' semantics to non-total orders leads to unintuitive judgements.

  $\models \varphi \diamond \rightarrow \psi$  iff:

$\exists \text{white circle} : \text{white circle} \models \varphi \wedge \forall \text{grey circle} : \text{grey circle} \models \varphi \Rightarrow \exists \text{green circle} : \text{green circle} \leq_{far} \text{grey circle} \wedge \text{green circle} \models \varphi$





# Fixing Lewis' Semantics

Non-total similarity relations

The semantics of  $\Box \rightarrow$  is too weak and of  $\Diamond \rightarrow$  too strong to capture the intended meaning on non-total relations.

We introduce operators with an additional level of quantification:

$\Box \bullet \rightarrow$  'Universal Would'

"If [...], **under all circumstances**,  $\psi$  would have been true as well."

$\Diamond \bullet \rightarrow$

'Existential Might'

"If [...], **under some circumstance**,  $\psi$  might have been true, too."



# Semantics of 'Universal Would'

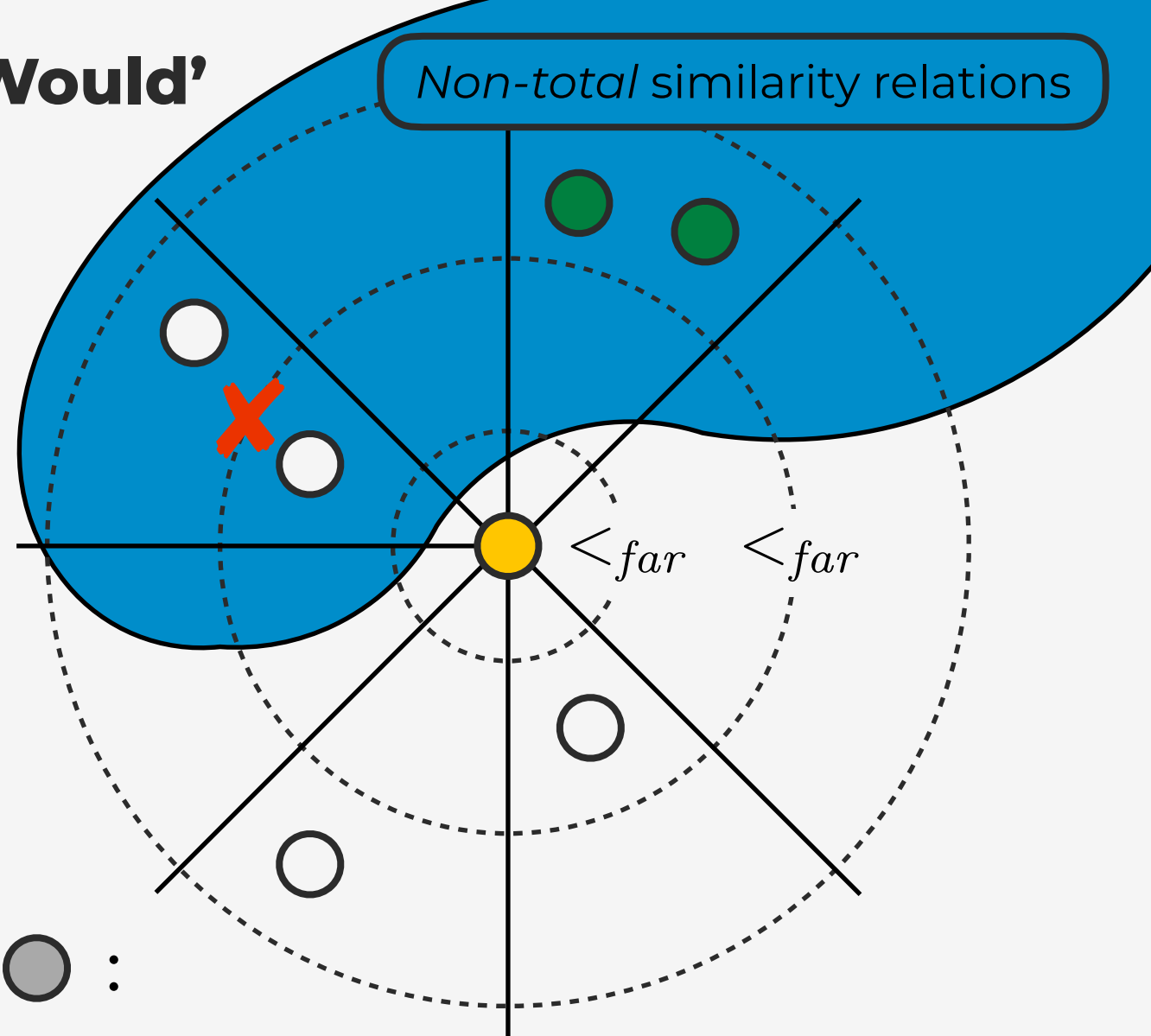
We quantify universally over  $\varphi$ -worlds and require a closer threshold world  $\bullet$  for all of them.

$\bullet \models \varphi \square \rightarrow \psi$  iff:

$\forall \circ : \circ \models \varphi \Rightarrow \exists \bullet :$

$\bullet \leq_{far} \circ \wedge \bullet \models \varphi \wedge \forall \circ :$

$\circ \leq_{far} \bullet \Rightarrow (\circ \models \varphi \rightarrow \psi)$







# Semantics of 'Universal Would'

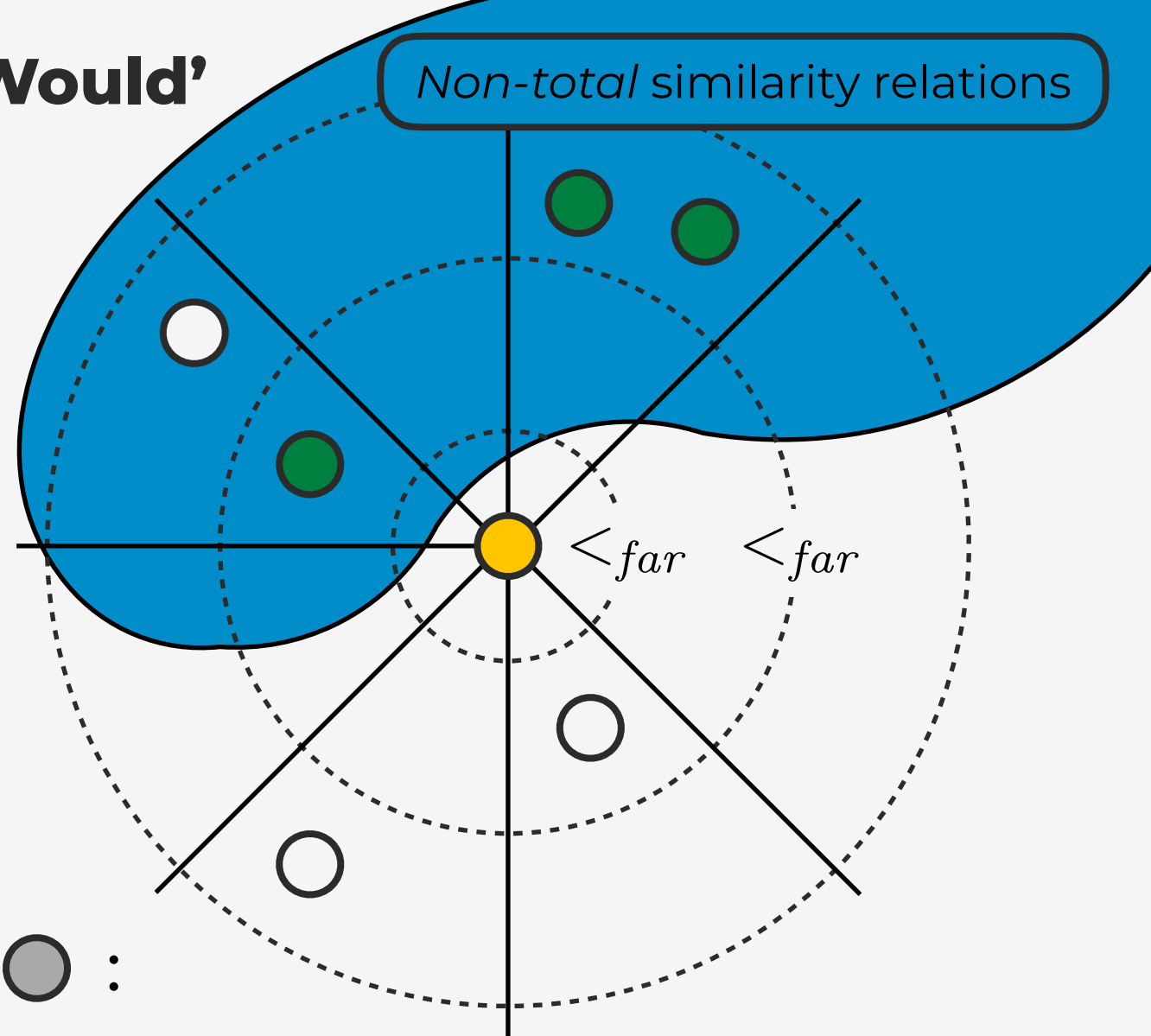
We quantify universally over  $\varphi$ -worlds and require a closer threshold world  $\bullet$  for all of them.

$\bullet$   $\models \varphi \square \rightarrow \psi$  iff:

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$\bullet \leq_{far} \circ \wedge \bullet \models \varphi \wedge \forall \circ :$

$\circ \leq_{far} \bullet \Rightarrow (\circ \models \varphi \rightarrow \psi)$





# Semantics of 'Existential Might'

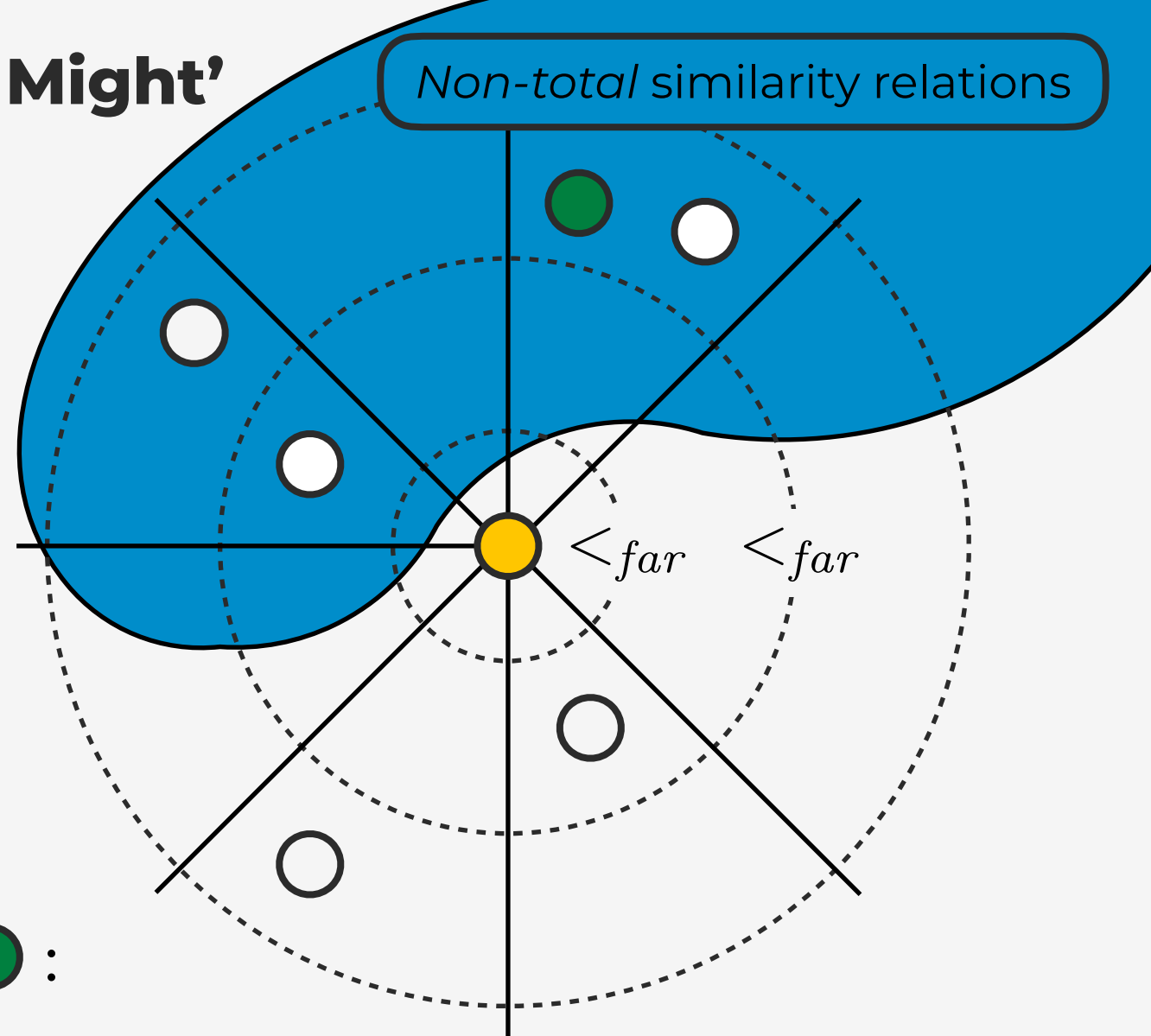
We quantify existentially over  $\varphi$ -worlds and require ever-closer worlds  $\bullet$  for one world only.

$\bullet$   $\models \varphi \blacklozenge \rightarrow \psi$  iff:

$\exists \circ : \circ \models \varphi \wedge \forall \bullet :$

$\bullet \leq_{far} \circ \wedge \bullet \models \varphi \Rightarrow \exists \bullet :$

$\bullet \leq_{far} \circ \wedge \bullet \models \varphi$

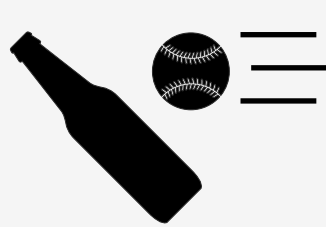




# Minimality

Minimal counterfactuals

For causality, causes are counterfactuals that describe the *minimal changes necessary* to avoid the effect.



Ball is thrown and bottle breaks.

< *far*



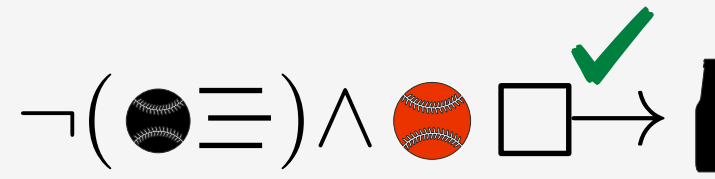
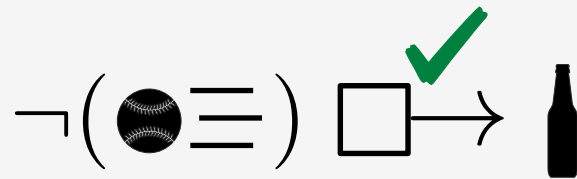
Ball is not thrown.



< *far*



Ball is not thrown and is red.



Antecedent is not *minimal*.



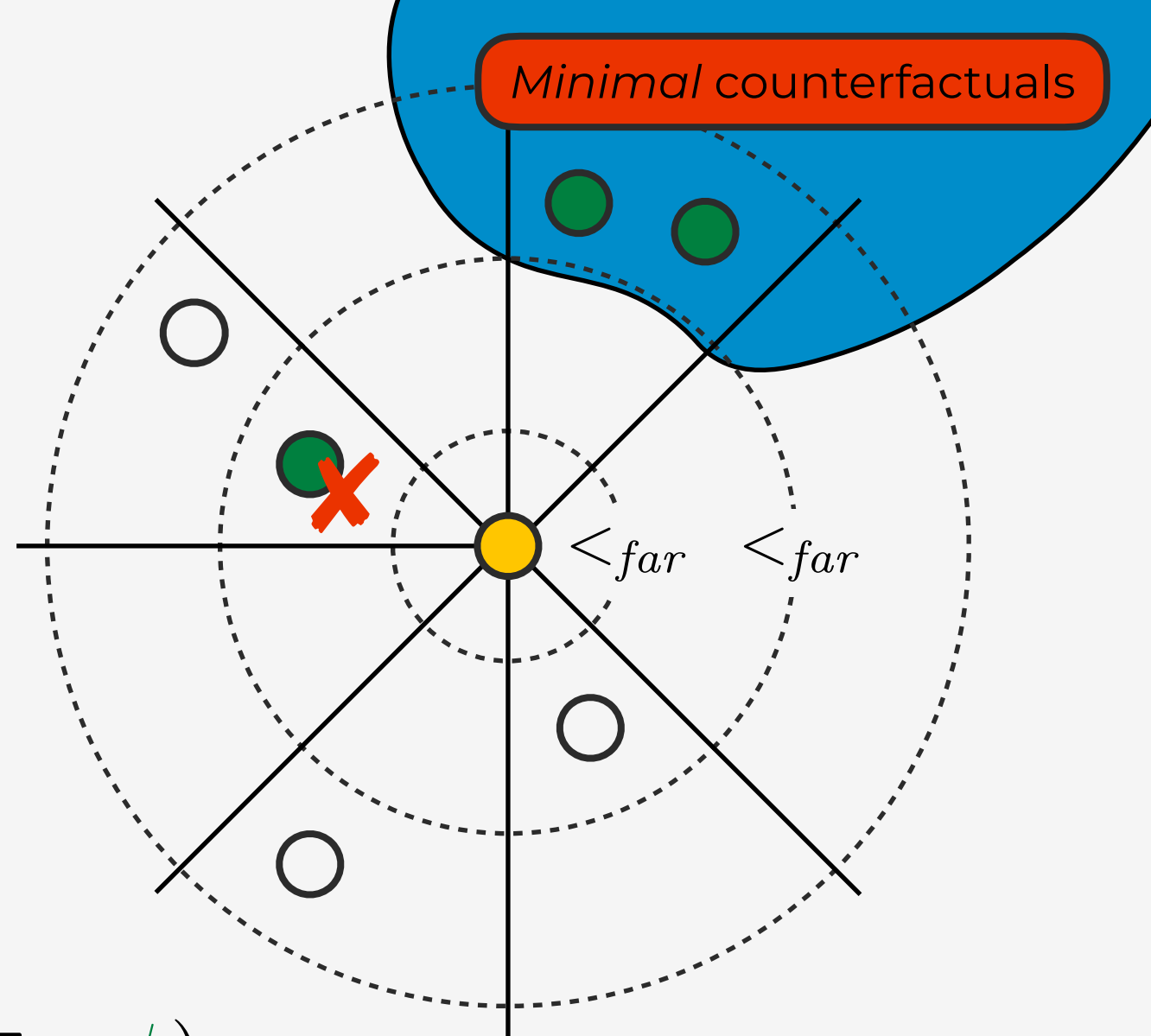
# Minimal Counterfactuals

A counterfactual is minimal if its antecedent describes the largest set that qualifies, e.g.,

$$\text{Yellow Circle} \models \varphi \square \rightarrow_{min} \psi \text{ iff:}$$

$$\text{Yellow Circle} \models \varphi \square \rightarrow \psi \wedge$$

$$\neg \exists \theta : \theta \supset \varphi \wedge (\text{Yellow Circle} \models \theta \square \rightarrow \psi)$$





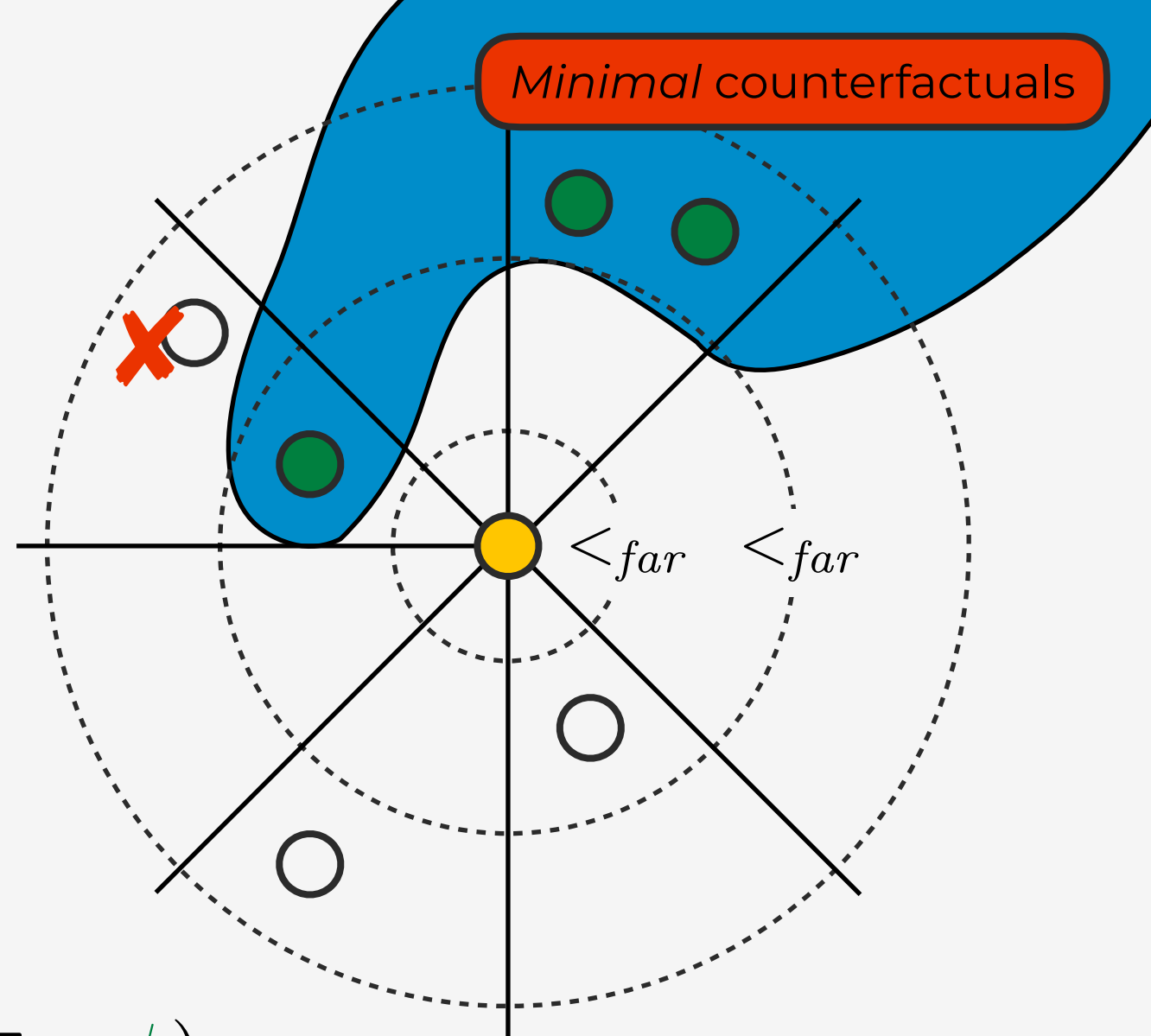
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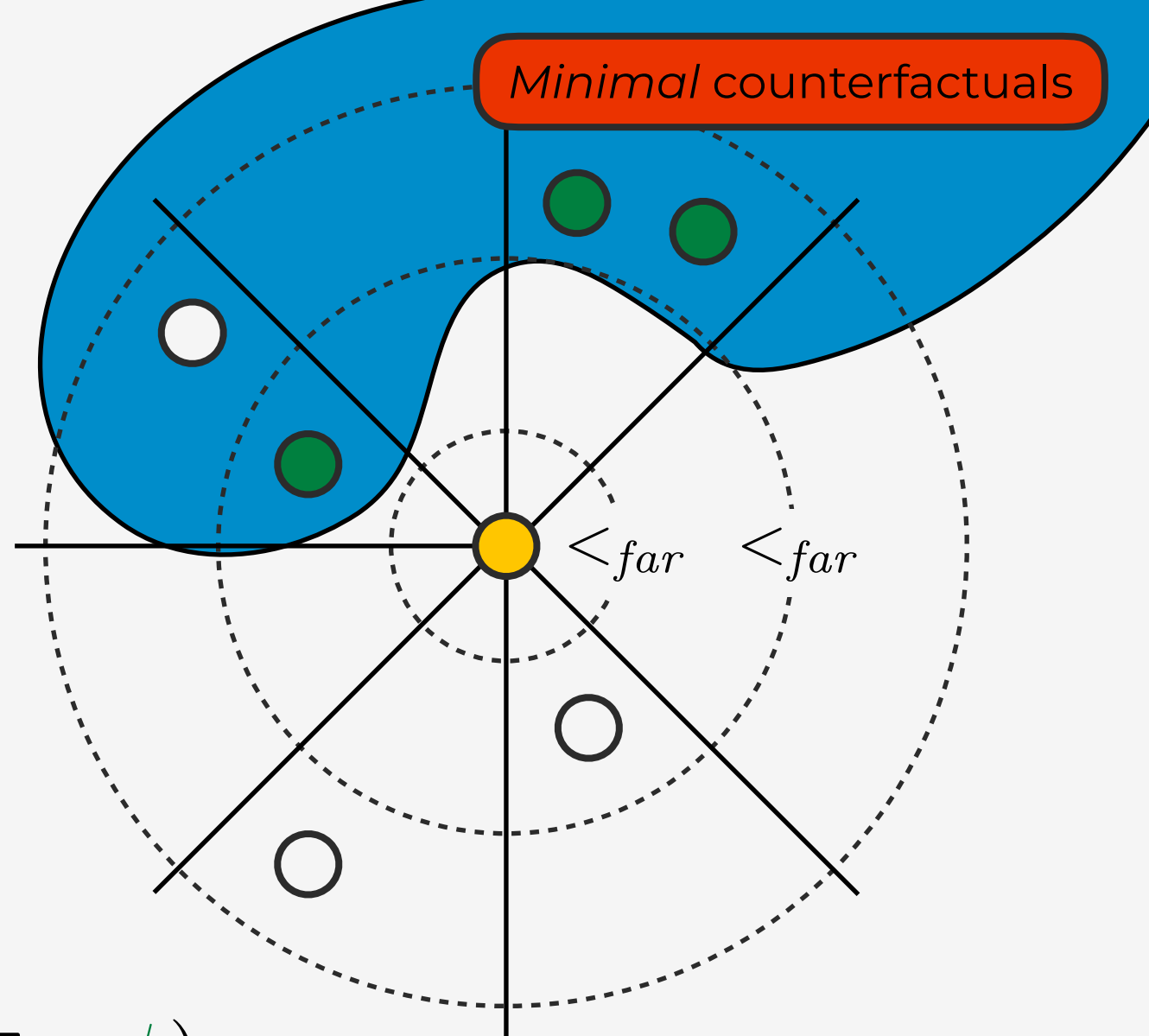
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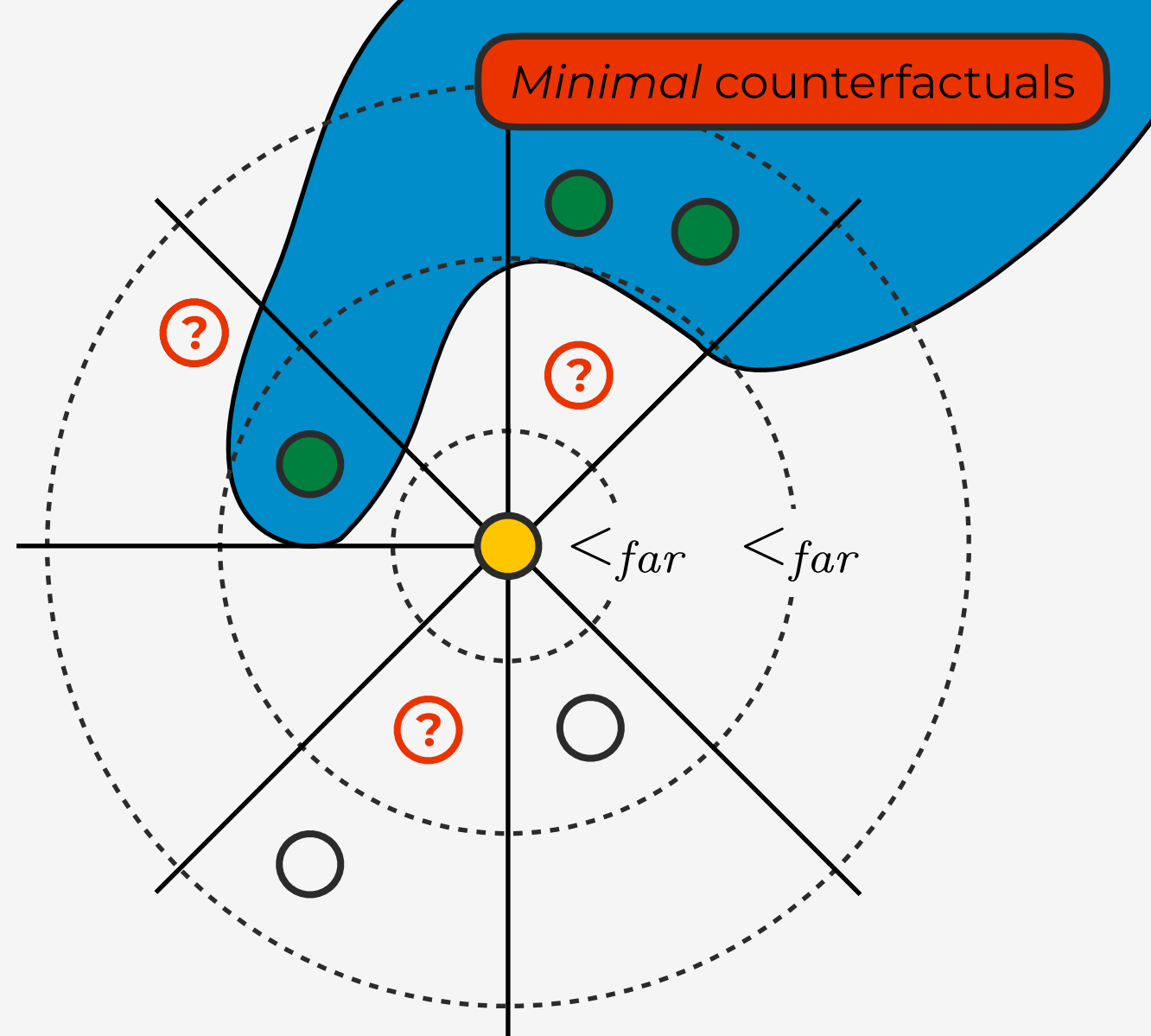


# Minimal Counterfactuals

The SO quantification can be avoided by searching for worlds  $\textcircled{?}$  that could be added to  $\varphi$ .

$\textcircled{\bullet} \models \varphi \quad \square \xrightarrow{\min} \psi$  iff:

Large FO-formula ( $\forall/\exists \textcircled{\bullet}$ )  
*See our paper!*





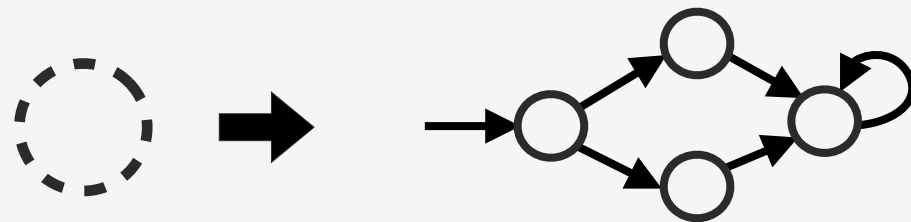
# Moving to Temporal Logics

$\varphi \square \mapsto \psi \quad \rightarrow \quad (\textit{straight}) \square \mapsto (\mathbb{F} \textit{ goal})$

$\bullet \quad \rightarrow \quad \textit{lasso trace, e.g. } \pi = \{a, b\}\{\}\{\{a, b\}\}^\omega$

$\circ \quad \rightarrow \quad \textit{infinite traces, e.g. } \pi' = \{a, b\}\{a, b\} \dots$

$\sigma \leq_{far}(\pi) \rho \quad \rightarrow \quad \textit{QPTL formula, e.g. } \bigwedge_{a \in AP} (a_\sigma \not\mapsto a_\pi) \rightarrow (a_\rho \not\mapsto a_\pi)$



$\forall/\exists \circ \quad \rightarrow \quad \forall/\exists \pi \quad (\textit{hyperproperty})$





# Satisfiability

Counterfactual TL

$$(\mathbf{X} a \square \rightarrow_{min} \mathbf{G} b) \wedge \dots$$

Counterfactuals Modulo QPTL



$$\forall/\exists \bigcirc$$

FO formula



$$\forall/\exists \pi$$

Prenex HyperQPTL



HyperQPTL Model Checking  
[Beutner&Finkbeiner, LPAR '23]

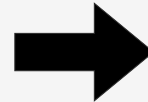


# Trace Checking

Counterfactual TL

$$(\mathbf{X} a \square \rightarrow_{min} \mathbf{G} b) \wedge \dots$$

Counterfactuals Modulo QPTL



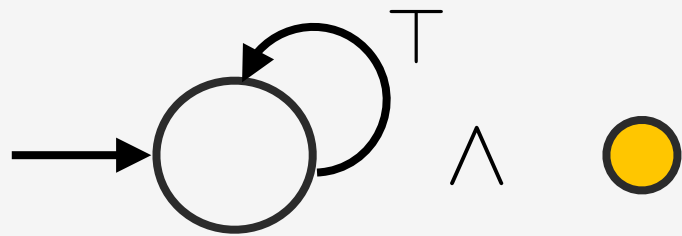
$$\forall/\exists \bigcirc$$

FO formula



$$\forall/\exists \pi$$

Prenex HyperQPTL



HyperQPTL Model Checking  
[Beutner&Finkbeiner, LPAR '23]



# Conclusion



Uniform specification language for counterfactual reasoning in, e.g., causality, fairness etc.



Automatic decision procedures for the resulting theory modulo QPTL.



System-Level Counterfactuals  
@*LPAR*: Counterfactuals Modulo Theories?