

Automata, Games, and Verification

Note: Due to the public holiday on the 2nd of June, there will be no tutorial next week. The solution of problem 1 will be presented in the lecture on the 31st of May. The second problem will be discussed in the tutorial on the 9th of June.

1. S1S and LTL

Decide for each of the languages over $2^{\{p,q\}}$ described below if they can be defined in S1S and/or LTL. Justify your answer in each case by either providing a formula or an argument why the language is not definable.

- $L_1 = \{\alpha \mid p \in \alpha(0), p \notin \alpha(i) \text{ for all } i \geq 1\}$;
- $L_2 = \{\alpha \mid p \in \alpha(i) \text{ for exactly two different } i \in \omega\}$;
- $L_3 = \{\alpha \mid |\{i \in \omega \mid p \in \alpha(i)\}| \text{ is finite and even}\}$;
- $L_4 = \{\alpha \mid |\{i \in \omega \mid p \in \alpha(i)\}| \text{ and } |\{i \in \omega \mid q \in \alpha(i)\}| \text{ are finite and equal}\}$.

2. S1S and LTL (tutorial A: group G03, tutorial B: group G06)

Let $L \subseteq (2^{\text{AP}})^\omega$ be an LTL-definable language and let $\text{AP}' \subsetneq \text{AP}$ be a strict subset of AP. Prove or give a counter example to the following statements:

- The (weak) projection $L_w = \{\sigma' \in (2^{\text{AP}'})^\omega \mid \exists \sigma \in L \forall i \in \omega. \sigma'(i) = \sigma(i) \cap \text{AP}'\}$ of L is LTL-definable.
- The (weak) projection L_w of L is S1S-definable.
- The strong projection $L_s = \{\sigma' \in (2^{\text{AP}'})^\omega \mid \forall \sigma \in (2^{\text{AP}})^\omega. (\forall i \in \omega. \sigma'(i) = \sigma(i) \cap \text{AP}') \rightarrow \sigma \in L\}$ of L is LTL-definable.
- The strong projection L_s of L is S1S-definable.