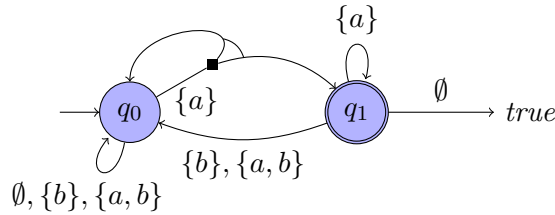


**Automata, Games, and Verification**

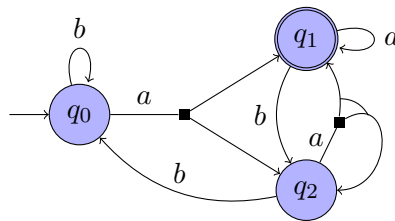
Please send a mail to [agv15@react.uni-saarland.de](mailto:agv15@react.uni-saarland.de) if you can't make it to the discussion session.

1. Consider the LTL formula  $\varphi = \square \diamond (\neg(a \mathcal{U} b))$  over the alphabet  $2^{\{a,b\}}$  and the alternating Büchi automaton  $\mathcal{A}$  given below. Does  $\mathcal{L}(\mathcal{A}) = \mathcal{L}(\varphi)$  hold?



- Yes       No

2. Consider the alternating Büchi automaton  $\mathcal{A}$  below. Which of the following are true?



- For all  $\alpha \in \mathcal{L}(\mathcal{A})$  there is a run tree on  $\alpha$  that is not memoryless.
  - For all  $\alpha \notin \mathcal{L}(\mathcal{A})$  there is a run tree on  $\alpha$  that is not memoryless.
  - There is some  $\alpha \in \Sigma^\omega$  such that all run trees on  $\alpha$  are memoryless.
  - There is some  $\alpha \in \Sigma^\omega$  such that all run trees on  $\alpha$  are not memoryless.
3. Let  $\mathcal{A}$  be some alternating Büchi automaton. Can there be words  $\alpha \in \mathcal{L}(\mathcal{A})$  that have an accepting run tree on  $\mathcal{A}$ , but no accepting run DAG?
- True       False
4. There is an alternating Büchi automaton  $\mathcal{A}$  and some word  $\alpha$  accepted by  $\mathcal{A}$  via some memoryless run  $r$ , whose corresponding run DAG is finite.
- True       False
5. Is there some accepting condition  $Acc$  and some alternating automaton  $\mathcal{A} = (\Sigma, Q, q_0, \delta, Acc)$  such that  $\mathcal{L}(\mathcal{A})$  is non-empty and no word  $\alpha$  accepted by  $\mathcal{A}$  has a memoryless run tree?
- Yes       No
6. There is an alternating Büchi automaton  $\mathcal{A}$  and some word  $\alpha$  such that there are infinitely many accepting run DAGs of  $\mathcal{A}$  on  $\alpha$ .
- True       False