Synthesis of Asynchronous Systems

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Motivation

- Why are asynchronous systems important?
- Natural way to model important problems
- Example:
 - distributed software modules
 - processes running at different speeds

Asynchronisity

- Programs take inputs from the environment
- Those inputs are invisible to the program until it is scheduled

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Output of the program may vary according to input

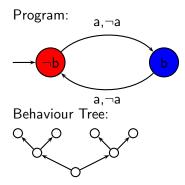
$$\underbrace{env} \xrightarrow{a} program} \xrightarrow{b}$$

Asynchronicity

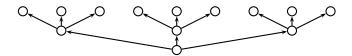
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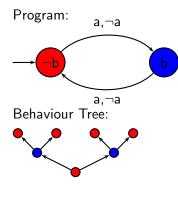
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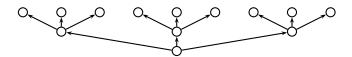
Computation Tree:



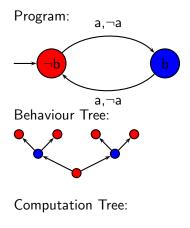
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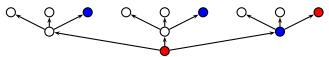


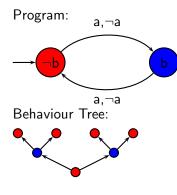
Computation Tree:



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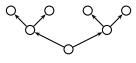
Computation Tree:



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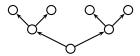
Synchronous Setting

Behaviour Tree:



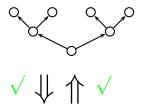
 $\checkmark \Downarrow$

Computation Tree:

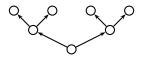


Synchronous Setting

Behaviour Tree:



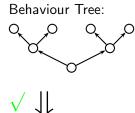
Computation Tree:



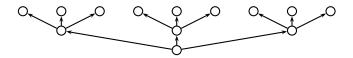
Computation tree and behaviour tree are essentially the same

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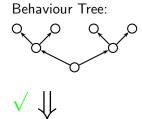
Behavour Tree to Computation Tree



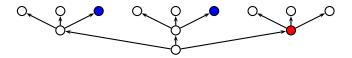
Computation Tree:



Computation Tree to Behaviour Tree?

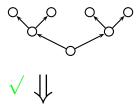


Computation Tree:

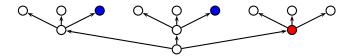


Computation Tree to Behaviour Tree?

Behaviour Tree:



Computation Tree:

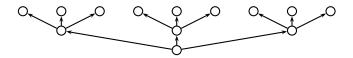


No corresponding behaviour tree

Computation Tree to Behaviour Tree? "No"

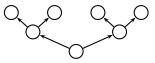
Behaviour Tree: $\checkmark \downarrow \uparrow \times$

Computation Tree:

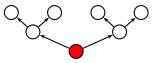


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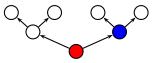
Computation Tree:

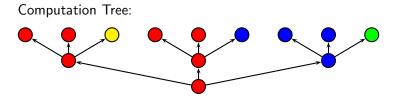


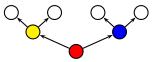
Computation Tree:

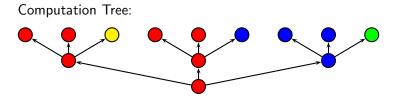


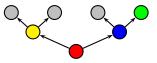
Computation Tree:











Deterministic Safety Tree Automata (DSTA)

- Acceptance game on CT
- Acceptance game on BT
- Automata Transformation
- Extensions

Automata Transformation

•
$$A_{CT} = (2^{I} \times 2^{S} \times 2^{O}, Q^{dir}, q_0, \delta)$$

- δ is a partial function
- $\blacktriangleright \ \delta(q, i, s, o) = (q_1, \textit{false}, \textit{env}) \land (q_2, \textit{true}, \textit{env}) \land (q_3, i, \textit{prog})$

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$$\bullet A_{BT} = (Qx2^{I}x2^{S}x2^{O}, dir, q_0, \delta')$$

- δ' is a partial function
- $\delta'(q, i, s, o) =$ $(q_1, false, env, \varepsilon) \land (q_2, true, env, \varepsilon) \land (q_3, i, prog, i)$

Getting Rid of ε Transitions

- Fixed Point Expansion
 - ► $\delta'(q, i, s, o) =$ $(q_1, false, env, o) \land (q_2, true, env, o) \land (q_3, i, prog, i)$

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Overview

► DSTA
$$A^{CT}$$
 \Rightarrow U ε STA A^{BT} \Rightarrow USTA A^{BT} \Rightarrow NET \Rightarrow Prog

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Extensions

- Universal Safety Tree Automata
 - ▶ we get it for free
- Universal Co-Buchi Automata
 - search for ε cycles on rejecting states
 - search for chains with infinitely many rejecting states
- LTL formulas could be represented as Universal Co-Buchi Automata

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Thanks For Listening

Questions?

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