



Exercise 10.1 - Recap (15)

1. Solving Muller games is in $\text{NP} \cap \text{Co-NP}$ if \mathcal{F} is encoded

- by a circuit?
 by a coloring function?
 by a tree?
 by an important subset?
 by a boolean formula?

2. In which of the following games is W_0 a trap for Player 1?

- Reachability
 Request Response
 Parity
 Safety
 Büchi
 Challenging Reachability
 Weak Muller

3. In a game with n vertices, which of the following winning conditions has the largest memory requirements for Player 0? (for sufficiently large n)

- Weak Muller
 Weak Parity
 Büchi
 Muller
 Parity
 Challenging Reachability
 Safety

4. What is the lower bound on the size of a winning strat. for Pl. 0 in weak Muller games?

- $\mathcal{O}(\log |\mathcal{F}|)$
 $\mathcal{O}(|V| \cdot |\mathcal{F}|)$
 $\mathcal{O}(2^{2^{|V|}})$
 $2^{\mathcal{O}(|V|)}$
 $\mathcal{O}(|V|^2)$
 $\mathcal{O}(\log^*(|V|))$
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5. Which of the following winning conditions can be described by a parity condition?

- SAFE
 REQRES
 WMULLER
 BÜCHI
 MULLER(\mathcal{F}) with \mathcal{F} doubly union-closed
 REACH
 CHREACH
 MULLER

6. Which of the following games are known to be solvable in polynomial time?

- Pushdown
 Bipartite Parity
 Büchi
 Solitary Parity
 Weak Parity
 Challenging Reachability
 Parity

7. Which of the following statements hold?

- $\text{BÜCHI}(\mathcal{F}) \in \Sigma_2$
 $\text{PARITY}(\Omega) \in \Sigma_2$
 $\text{WMULLER}(\mathcal{F}) \in \Sigma_2 \cap \Pi_2$
 $\text{SAFE}(S) \in \Sigma_2 \cap \Pi_2$
 $\text{REQRES}((Q_j, P_j)_{j \in [k]}) \in \Sigma_3$
 $\text{BÜCHI}(\mathcal{F}) \in \Pi_2$
 $\text{MULLER}(\mathcal{F}) \in \Pi_2$
 $\text{REACH}(R) \in \Pi_1$

8. Which of the following winning conditions are prefix independent?

- WMULLER
 COBÜCHI
 BÜCHI
 REQRES
 MULLER
 PARITY
 SAFE

9. For which of the following games has Player 1 positional winning strategies?

- Muller
 Büchi
 Reachability
 co-Büchi
 Weak Parity
 Safety
 Parity
 Challenging Reachability

10. In which of the following games can Player 0 win with a uniform strategy?

- Muller
 Reachability
 Challenging Reachability
 co-Büchi
 Weak Parity
 Request-Response
 Parity

11. Let $\mathcal{G} = (\mathcal{A}, \text{PARITY}(\Omega: V \rightarrow [k]))$ be a parity game with $\text{Par}(k) = 0$. How large is $|\text{Sh}(\mathcal{G})|$?

- $|V|$
 $2^{|V|}$
 $|V|^2$
 $1 + \prod_{c \in \Omega(V), \text{Par}(c)=1} |\Omega^{-1}(c)|$
 $\frac{|V|!}{2} + 1$
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 $|V| \cdot |E|$

12. Which of the following games are self-dual?

- co-Büchi
 Request-Response
 Reachability
 Parity
 Challenging Reachability
 Weak Muller
 Weak Parity

13. In which of the following games may Player 0 need memory?

- Büchi
 Weak Parity
 Safety
 Weak Muller
 Parity
 Request-Response
 Muller

14. Using game reductions, can you reduce

- co-Büchi to Reachability?
 Büchi to Safety?
 Safety to Reachability?
 Parity to Muller?
 Muller to Büchi?
 Request-Response to Parity?
 Büchi to Parity?

15. Which of the following games are determined?

- Muller
 Safety
 Request-Response
 Parity
 Büchi
 Reachability
 co-Büchi

Exercise 10.2 - Zielonka Trees, Rabin, Streett

(10 Bonus Points)

Given a family $\mathcal{F} \subseteq 2^V$ of subsets of a finite set V , we define its Zielonka tree $\mathcal{Z}(\mathcal{F})$ recursively as follows:

- The root of $\mathcal{Z}(\mathcal{F})$ is labeled by the set of all vertices.
- Children of a node labeled with $F \in \mathcal{F}$ are the \subseteq -maximal subsets $F' \subseteq F$ with $F' \notin \mathcal{F}$.
- Children of a node labeled with $F \notin \mathcal{F}$ are the \subseteq -maximal subsets $F' \subseteq F$ with $F' \in \mathcal{F}$.

We already had an example of such a tree on page 39 of the lecture notes. We say that a vertex v of $\mathcal{Z}(\mathcal{F})$ is a Player 0 vertex if and only if its label is in \mathcal{F} .

Given a family $(Q_j, P_j)_{j \in [k]}$ of subsets $Q_j, P_j \subseteq V$ with $k \in \mathbb{N}$ we define the Rabin winning condition by

$$\text{RABIN}((Q_j, P_j)_{j \in [k]}) = \{ \rho \in V^\omega \mid \exists j \in [k]. \text{Inf}(\rho) \cap Q_j \neq \emptyset \wedge \text{Inf}(\rho) \cap P_j = \emptyset \}$$

and the Streett winning condition by

$$\text{STREETT}((Q_j, P_j)_{j \in [k]}) = \{ \rho \in V^\omega \mid \forall j \in [k]. \text{Inf}(\rho) \cap Q_j \neq \emptyset \Rightarrow \text{Inf}(\rho) \cap P_j \neq \emptyset \}$$

Given an arena $\mathcal{A} = (V, V_0, V_1, E)$ we then call the games $\mathcal{G}_r = (\mathcal{A}, \text{RABIN}((Q_j, P_j)_{j \in [k]}))$ and $\mathcal{G}_s = (\mathcal{A}, \text{STREETT}((Q_j, P_j)_{j \in [k]}))$ a Rabin game and a Street game, respectively.

Prove the following statements:

- a) For every family $(Q_j, P_j)_{j \in [k]}$ with $j \in \mathbb{N}$ and $Q_j, P_j \subseteq V$ holds that

$$\text{RABIN}((Q_j, P_j)_{j \in [k]}) = V^\omega \setminus \text{STREETT}((Q_j, P_j)_{j \in [k]}).$$

- b) For every coloring function $\Omega: V \rightarrow \mathbb{N}$ there exists a family $(Q_j, P_j)_{j \in [k]}$ with $j \in \mathbb{N}$ and $Q_j, P_j \subseteq V$ such that $\text{PARITY}(\Omega) = \text{RABIN}((Q_j, P_j)_{j \in [k]})$.
- c) For every coloring function $\Omega: V \rightarrow \mathbb{N}$ there exists a family $(Q_j, P_j)_{j \in [k]}$ with $j \in \mathbb{N}$ and $Q_j, P_j \subseteq V$ such that $\text{PARITY}(\Omega) = \text{STREETT}((Q_j, P_j)_{j \in [k]})$.
- d) For every family $(Q_j, P_j)_{j \in [k]}$ with $j \in \mathbb{N}$ and $Q_j, P_j \subseteq V$ there is a set $\mathcal{F} \subseteq 2^V$ such that $\text{RABIN}((Q_j, P_j)_{j \in [k]}) = \text{MULLER}(\mathcal{F})$.
- e) For every family $(Q_j, P_j)_{j \in [k]}$ with $j \in \mathbb{N}$ and $Q_j, P_j \subseteq V$ there is a set $\mathcal{F} \subseteq 2^V$ such that $\text{STREETT}((Q_j, P_j)_{j \in [k]}) = \text{MULLER}(\mathcal{F})$.
- f) For every set $\mathcal{F} \subseteq 2^V$ holds: every Player 0 vertex of $\mathcal{Z}(\mathcal{F})$ has at most one successor if and only if $\text{MULLER}(\mathcal{F}) = \text{RABIN}((Q_j, P_j)_{j \in [k]})$ for some family $(Q_j, P_j)_{j \in [k]}$ with $j \in \mathbb{N}$ and $Q_j, P_j \subseteq V$.
- g) For every set $\mathcal{F} \subseteq 2^V$ holds: every Player 1 vertex of $\mathcal{Z}(\mathcal{F})$ has at most one successor if and only if $\text{MULLER}(\mathcal{F}) = \text{STREETT}((Q_j, P_j)_{j \in [k]})$ for some family $(Q_j, P_j)_{j \in [k]}$ with $j \in \mathbb{N}$ and $Q_j, P_j \subseteq V$.
- h) For every set $\mathcal{F} \subseteq 2^V$ holds: every vertex of $\mathcal{Z}(\mathcal{F})$ has at most one successor if and only if $\text{MULLER}(\mathcal{F}) = \text{PARITY}(\Omega)$ for some coloring function $\Omega: V \rightarrow \mathbb{N}$.
- i) Let $\mathcal{Z}(\mathcal{F})$ be the Zielonka tree for some $\mathcal{F} \subseteq 2^V$ such that there is a Player i vertex of $\mathcal{Z}(\mathcal{F})$ which has two successors. Then there is a Muller game $\mathcal{G} = (\mathcal{A}, \text{MULLER}(\mathcal{F}))$ with vertex set V where Player i has a winning strategy from some $v \in V$, but no positional one.
- j) For every \mathcal{F}_n with $n \in \mathbb{N}^+$ defined as in the game DJW_n we have that $\mathcal{Z}(\mathcal{F}_n)$ has at least $n!$ many leaves.